Schlecht, Sebastian J; Habets, Emanuël A P

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Modal Decomposition of Feedback Delay Networks

Sebastian J. Schlecht and Emanuël A. P. Habets, Senior Member, IEEE

Abstract—Feedback delay networks (FDNs) belong to a general class of recursive filters which are widely used in sound synthesis and physical modeling applications. We present a numerical technique to compute the modal decomposition of the FDN transfer function. The proposed pole finding algorithm is based on the Ehrlich-Aberth iteration for matrix polynomials and has improved computational performance of up to three orders of magnitude compared to a scalar polynomial root finder. The computational performance is further improved by bounds on the pole location and an approximate iteration step. We demonstrate how explicit knowledge of the FDN’s modal behavior facilitates analysis and improvements for artificial reverberation. The statistical distribution of mode frequency and residue magnitudes demonstrate that relatively few modes contribute a large portion of impulse response energy.

Index Terms—Feedback Delay Network, Modal Synthesis, Artificial Reverberation, Matrix Polynomial, Ehrlich-Aberth Iteration

I. INTRODUCTION

A feedback delay network (FDN) consists of a set of an \( N \) delay lines with lengths \( m \) which are interconnected via a feedback matrix \( A \) (see Fig. 1). FDNs arise in many physical modeling applications where geometrically distributed components are approximated by time delays, e.g., strings [1], plates and membranes [2], springs [3], and air volume [4, 5]. The interest in FDNs is fueled by the highly efficient implementation of delays in the time-domain, e.g., with circular buffers resulting in a constant time complexity \( O(1) \) independent of its length. Therefore, the computational complexity of the FDN scales with the number of delay lines and not the system order. FDNs are a popular choice for artificial reverberation applications particularly because of the favorable relation between FDN size and system order [6–9].

In this work, we present a modal decomposition technique for FDNs. The modal decomposition of a system is an equivalent representation as the sum of complex one-pole resonators, so-called modes. The time-domain signal of such a resonator with pole \( \lambda_i \) and residue \( \rho_i \) is

\[
h_i(n) = |\rho_i| |\lambda_i|^n e^{i(n\angle \lambda_i + \angle \rho_i)},
\]

where \( \angle \) indicates the argument of a complex number in radians, \( |\cdot| \) is the magnitude, \( i = \sqrt{-1} \) and \( n \) indicates the discrete time index. Each individual resonating mode is governed by four parameters: mode frequency \( \angle \lambda_i \), decay rate \( \angle \rho_i \), initial phase \( \angle \rho_i \) and initial amplitude \( |\rho_i| \) (see Fig. 1).

With modal decomposition, we aim to uncover the specific parameters of each mode. The time-domain impulse response of the FDN

\[
h(n) = \sum_{i=1}^{\aleph} h_i(n)
\]

is the sum of the complex modes \( h_i(n) \), where \( \aleph \) is the system order. In sound synthesis applications for instance, the human auditory system can recognize the spectral quality composed of the individual modes and this representation is therefore termed additive or modal synthesis [10]. Modal analysis of recursive systems is applied in various system modeling applications, ranging from acoustics and digital filter design to mechanical modeling [11]. A particularly challenging application for modal decomposition is room acoustics, where even rooms of medium sizes exhibit millions of modes [12]. Only for simple room geometries, an analytic expression for the system poles and residues can be stated [13]. System poles may also be recovered from the impulse response by various techniques such as an autoregressive moving-average [14], Bayesian inference [15] and all-pole modeling [16]. Whereas these techniques may be able to successfully compute partial solutions or compute the solution for specific configurations, the computation of the entire set of modes is in general challenging.

In the remainder of this paper, we present a numerically stable and computationally efficient method to compute the modal decomposition for large system order \( \aleph \) and modest-sized \( N \). The proposed pole finding algorithm is based on the Ehrlich-Aberth iteration for matrix polynomials and has improved computational performance of up to three orders
of magnitude compared to a scalar polynomial root finder. In Section II, we state the precise problem and review the prior art. In Section III, we derive bounds on the pole locations and improved initial estimates. In Section IV, we propose a computationally efficient method to estimate the pole locations. In Section V, we evaluate the performance of the proposed method. In Section VI, we apply modal decomposition to analyze the effects of attenuation filters and to study the statistical distributions of mode frequencies and residue magnitudes. Further, we demonstrate how explicit knowledge of the FDN’s modal behavior facilitates analysis and improvements for artificial reverberation.

II. PROBLEM STATEMENT AND PRIOR ART

In the following, we give the precise problem statement of this work and the prior art.

A. Problem Statement

For a single input and single output, the time-domain recursion of an FDN with \( N \) delay lines is given by

\[
y(n) = c^T s(n) + d x(n),
\]

\[
s(n + m) = A s(n) + b x(n),
\]

where \( n \) is the time index, \( \cdot^T \) denotes the transpose operation and \( A \in \mathbb{C}^{N \times N} \), \( b, c, s(n) \in \mathbb{C}^{N \times 1} \), \( x(n), y(n), d \in \mathbb{C} \) [17]. The state vector is defined as \( s(n + m) = [s_1(n + m_1), \ldots, s_N(n + m_N)] \). We write \( N \)-FDN to denote an FDN of size \( N \). The transfer function of an FDN is

\[
H(z) = c^T [D_m(z)^{-1} - A]^{-1} b + d,
\]

where \( D_m(z) = \text{diag}(z^{-m_1}, \ldots, z^{-m_N}) \). The system order is given by \( \mathcal{N} = \sum_{i=1}^{N} m_i \) [17]. For commonly used delays \( m \), the system order is much larger than the FDN size, i.e.,

\[
\mathcal{N} \gg N.
\]

The modal decomposition of the FDN, i.e., the partial fraction decomposition (PFD) of the transfer function (4), is

\[
H(z) = d + \sum_{i=1}^{\mathcal{N}} \frac{\rho_i}{1 - \lambda_i z^{-1}},
\]

where \( \rho_i \) is the residue of the pole \( \lambda_i \). The time-domain representation of the sum in (6) is given in (2) as the sum of complex resonators. The objective of this work is to present an efficient numerical method to compute the modal decomposition (6) from the transfer function (4).

B. Direct Approach

We first review two standard methods for the modal decomposition [17, 18]. Let \( A \) be an invertible matrix, then \( \text{adj}(A) = \det(A)A^{-1} \) is the adjugate of the matrix \( A \) [19]. In the following, we denote

\[
P(z) = D_m(z)^{-1} - A.
\]

With (7), the transfer function (4) can be expressed as a rational polynomial

\[
H(z) = \frac{q_{m,A,b,c,d}(z)}{p_{m,A}(z)},
\]

where

\[
p_{m,A}(z) = \det(P(z)) \tag{9}
\]

and

\[
q_{m,A,b,c,d}(z) = d \det(P(z)) + c^T \text{adj}(P(z)) b. \tag{10}
\]

For brevity, we occasionally omit the parameters and write \( q(z) \) and \( p(z) \). The FDN system poles \( \lambda_i \), where \( 1 \leq i \leq \mathcal{N} \), are the roots of the generalized characteristic polynomial (GCP) \( p_{m,A}(z) \) in (9) such that they are fully characterized by the delays \( m \) and the feedback matrix \( A \).

For a moment, let us assume that all delays are single time steps, i.e., \( m = 1 \). The time-domain recursion in (3) reduces to the standard state-space description of a linear time-invariant (LTI) filter. The system poles \( \lambda_i \) are the eigenvalues of the feedback matrix \( A \) such that the modal decomposition (6) is easily computed with standard methods. However, for longer delays \( m \) such that (5) holds, the modal decomposition becomes more involved.

The GCP \( p_{m,A} \) can be expanded in a linearized fashion

\[
p_{m,A}(z) = \det(z I_{\mathcal{N}} - A), \tag{11}
\]

where \( I_{\mathcal{N}} \) is the identity matrix of size \( \mathcal{N} \) and \( A \in \mathbb{C}^{\mathcal{N} \times \mathcal{N}} \) such that the system poles are the eigenvalues of \( A \) [17]. Unfortunately, for large delays \( m \) this eigenvalue problem quickly becomes numerically intractable as demonstrated in Section V-B. Alternatively, the GCP can be expressed as a scalar polynomial

\[
p_{m,A}(z) = \sum_{i=0}^{\mathcal{N}} c_i z^i, \tag{12}
\]

where the coefficients \( c_i \) are derived from the principal minors of \( A \) [18]. The system poles are the roots of the scalar polynomial. Again, the polynomial degree increases with longer delays \( m \) and finding the roots of the polynomial becomes numerically intractable [20] (see Section V-B).

C. Ehrlich-Aberth Method

In the following, we review the Ehrlich-Aberth method for determining polynomial eigenvalues based on the polynomial matrix formulation of the FDN. While the method itself is prior art, we like to emphasize that the application to FDNs is novel.

1) Polynomial Matrix Formulation: It is a common heuristic in numerical computation that the inherent problem structure shall be preserved as much as possible throughout all computation steps to improve numerical performance. In contrast to Section II-B, we compute the system poles without expanding the problem. In fact, (9) is a polynomial eigenvalue problem of degree \( K = \max_m m \), i.e.,

\[
P(z) = \sum_{k=0}^{K} P_k z^k, \tag{13}
\]
where \( P_k \in \mathbb{C}^{N \times N} \) for \( 0 \leq k \leq K \). For a proper matrix polynomial \( P(z) \), i.e., \( \det(P) \neq 0 \), the number of roots is \( KN \) [21]. For FDNs, however, the matrix polynomial \( P(z) \) in (7) is improper because \( \det(P_K) = 0 \). In fact, if \( \det(A) \neq 0 \), the number of finite roots is \( \Re \) which is also the degree of the scalar polynomial in (9) [18].

In the following, we use the derivative of the polynomial \( p(z) = \det(P(z)) \). According to Jacobi’s formula [22], we have

\[
p'(z) = \frac{d}{dz} p(z) = \det(P(z)) \text{tr}(P(z)^{-1} P'(z))
\]

where \( P'(z) = \frac{dP(z)}{dz} \) and \( \text{tr}(X) \) denotes the trace of matrix \( X \). Stewart [23] showed that the adjugate of \( A \) can be well-conditioned even when \( A \) is ill-conditioned, and he shows how \( \text{adj}(A) \) can be computed in a numerically stable way from a rank revealing decomposition of \( A \) [22]. With the definitions of the polynomial eigenvalue problem introduced, we present the core iteration step.

2) Ehrlich-Aberth Iteration: The polynomial eigenvalue problem can be solved with the Ehrlich-Aberth Iteration (EAI) method, i.e., a combination of the Newton method and a deflation term which prevents that two eigenvalues converge to the same solution [21]. Let \( \lambda^{(0)} \subset \mathbb{C} \) be a set of initial estimates for the \( \Re \) roots of the polynomial \( p(z) \) and \( \lambda^{(j)} = \{ \lambda_1^{(j)}, \lambda_2^{(j)}, \ldots, \lambda_{\Re}^{(j)} \} \) be the \( j \)-th EAI iteration. We assume that the phase of the estimates \( \lambda^{(j)} \) are sorted in ascending order, i.e., \( \angle \lambda_1^{(j)} < \angle \lambda_i^{(j)} \) if \( i < l \). The EAI provides the sequence of estimates

\[
\lambda_i^{(j+1)} = \lambda_i^{(j)} - \Delta_i^{(j)}
\]

with the EAI step being

\[
\Delta_i^{(j)} = \frac{1}{\mathcal{N}(\lambda_i^{(j)})} - D_i(\lambda^{(j)})
\]

Using the identity in (14), the Newton correction term is

\[
\mathcal{N}(z) = \frac{p(z)}{p'(z)} = \frac{1}{\text{tr}(P(z)^{-1} P'(z))}
\]

and the deflation term is

\[
D_i(\lambda^{(j)}) = \sum_{l=1, l \neq i}^{\Re} \frac{1}{\lambda_i^{(j)} - \lambda_l^{(j)}}.
\]

The deflation term may be interpreted as a penalty term if two eigenvalues approach each other too closely and guarantees that the all eigenvalues reached are unique. The method, given here in the Jacobi version, is known to converge cubically for simple roots and linearly for multiple roots [21]. The Gauss-Seidel version of EAI [21], which updates the estimates as soon as they become available, may converge even slightly faster.

3) Stopping Criteria: The system poles \( \lambda_i \) are the roots of the polynomial \( p(z) \) in (9), i.e., \( \det(P(\lambda)) = 0 \). In other words, \( P(\lambda_i) \) is a singular matrix for all system poles. Thus, the natural stopping criteria is the reciprocal of the condition number \( \kappa(P(z)) \) being less than a prescribed tolerance \( \tau_1 \). This stopping condition is also computationally favorable as the condition number can be estimated highly efficiently [22]. However, for multiple eigenvalues this stopping condition may result in a premature halt [21].

An alternative stopping condition says that the computed correction is too tiny and would not change the significant digits of the current estimate

\[
\left| \Delta_i^{(j)} \right| \leq \tau_2 \left| \lambda_i^{(j)} \right|
\]

where \( \tau_2 \) is a small positive tolerance threshold [21]. In practice, good global convergence properties are observed; a theoretical analysis of global convergence, though, is still missing and constitutes an open problem. There is empirical evidence that the number of Newton iterations heavily depends on the choice of the initial estimates [21].

D. Residues

Once we have found the system poles, the residues of the modal decomposition (6) are computed by

\[
\rho_i = \frac{q(\lambda_i)}{p'(\lambda_i)},
\]

where we assume that all poles are unique. Similar, but more intricate solutions exists for non-unique poles [24]. The undriven residue, i.e., the system response without excitation, is

\[
\rho_i^n = \frac{1}{p'(\lambda_i)}.
\]

The undriven residue is a valuable intermediate step to analyze the mode initial amplitude independent from the input and output drives \( q(\lambda_i) \). Since, \( P(\lambda_i) \) is a singular matrix, the derivative of the GCP \( p'(\lambda_i) \) in (14) may only be computed by the adjugate formulation. Since \( \det(P(\lambda_i)) = 0 \), the input-output drives in (10) are

\[
q(\lambda_i) = \mathbf{c}^\top \text{adj}(P(\lambda_i)) \mathbf{b}.
\]

The difference between the driven and undriven residues may be expressed as a linear combination of the matrix entries of \( \text{adj}(P(\lambda_i)) \). Alternatively, the driven residues may also be computed by a least linear squares fit for the time-domain impulse response since the sum of complex resonators in (2) depends linearly on the residues [25].

III. Bounds on the Pole Location

For FDNs, the near-equidistribution of mode frequencies has been conjectured before [26] and the authors have given an analytical bound on the equidistribution based on Hayman’s theorem [18]. We give further experimental evidence for the near-equidistribution in Section V-B. Aberth [27] proposed to choose initial estimates placed along a circle centered at the origin of sufficiently large radius so that it contains all the
roots. In case the magnitude of the roots vary largely, multiple circles with suitable radii may be chosen instead [28].

In this section, we establish novel bounds for the pole locations of FDNs. Such bounds are essential for good initialization which can improve the performance of the algorithm. We conclude this section with a generalization to additional filtering in the delay lines and feedback matrix.

A. Frequency-Independent Bounds

With Rouche’s theorem, we can derive upper and lower bounds on the pole magnitudes for the FDN depending on the singular values $\sigma(A)$ of the feedback matrix (see Appendix A)

$$\min \frac{m}{\min \sigma(A)} \leq |\lambda_i| \leq \max \frac{m}{\max \sigma(A)}.$$  \hspace{1cm} (23)

Equation (23) is a generalization on the relation of eigenvalues and singular values as given in the Weyl-Horn Theorem [29] in the case of unit delays $m = 1$. The bound is tight for a diagonal feedback matrix where the minimum and maximum delays coincide with the minimum and maximum diagonal element, respectively. However, the bound may be arbitrarily loose. For instance, the maximum singular value of a triangular matrix $\max \sigma(A)$ may be arbitrarily large while all system poles lie on the unit circle [9]. For large delays $m$ however, (23) shows that the pole magnitudes tend to be close to the unit circle.

We can further derive from (23) that if $\max \sigma(A) \leq 1$ then all poles lie on the closed unit disk which is equivalent to the FDN being marginally stable if all poles on the unit circle are simple, i.e., have algebraic multiplicity equal to 1 [30]. In particular, if all singular values are 1, which is equivalent to $A$ being unitary, i.e., $A^H A = I$, all system poles lie on the unit circle regardless of the delays $m$. Such an FDN is called lossless, and represents an important special case [9].

B. Polynomial Feedback and Delay Matrices

Although, the focus of this work is on frequency-independent feedback matrices $A$, much of the development in this work is applicable to general polynomial matrices. Therefore, it is easy to include further filtering such as a frequency-dependent feedback matrix $A(z)$. There also exists a singular value decomposition for polynomial matrices $A(z)$ [31]. Alternatively, the delay lines are often extended with an attenuation filter $\alpha_i(z)$, i.e.,

$$D_m(z) \text{diag}(\alpha(z)) = \text{diag}(z^{-m_1} \alpha_1(z), \ldots , z^{-m_N} \alpha_N(z)).$$  \hspace{1cm} (24)

It is important to note that additional filters may increase the number of system poles. Further, if $P(z)$ is a rational polynomial, in other words consists of IIR filters, then the transfer function in (6) is no longer proper, i.e., the polynomial degree of the nominator is larger than the polynomial degree of the denominator [32]. Nonetheless, improper partial fraction decomposition can be solved with a delayed parallel form by separating the FIR and IIR part of the transfer function [25].

For a unitary feedback matrix $A$, large delays $m$, and, for attenuation filters in (24), it is possible to improve the pole magnitude bounds (23) to

$$\min \left( |\alpha(z) e^{j\lambda_i}| \right)^{1/m} \leq |\lambda_i| \leq \max \left( |\alpha(z) e^{j\lambda_i}| \right)^{1/m},$$  \hspace{1cm} (25)

where all vector operations are element-wise (see Appendix A).

C. Initial Estimates

For lossless FDNs, we place the initial estimates $\lambda^{(0)}$ uniformly on the unit circle. More precisely, we chose the roots of unity

$$\lambda^{(0)} = \exp \left( j2\pi \left[ \left. \frac{0}{M}, \frac{1}{M}, \ldots , \frac{M-1}{M} \right] \right).$$  \hspace{1cm} (26)

It is worthwhile to note that $\lambda^{(0)}$ is the solution of a particular FDN with a circular shift matrix

$$A = I_S = \begin{bmatrix} 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \\ 1 & 0 & 0 & \cdots & 0 \end{bmatrix}$$  \hspace{1cm} (27)

such that the GCP is

$$p_m(z) \lambda^{(0)} = z^M - 1.$$  \hspace{1cm} (28)

The shift matrix thus combines the FDN delays into a single long delay line. For non-lossless FDNs, we employ the bounds on the pole locations (23) and (25) for improved initial estimates by placing the initial poles along the maximum bounding curve.

IV. IMPROVED EHRLICH-ABERTH ITERATION

In the previous section, we have shown that for large delays $m$, the poles tend to be close to the unit circle with nearly equidistributed phases. Based on this observation, we propose in this section an improved iteration step.

A. Approximate Deflation

For a high system order $\mathcal{O}$, the computational complexity of the deflation term (18) may become excessive. We propose an approximate deflation (EAI) according to a maximum error tolerance $\tau_3$ for the resulting EAI step $\Delta_i^{(j)}$ in (16), i.e.,

$$|\Delta_i^{(j)} - \overline{\Delta}_i^{(j)}| \leq \tau_3.$$  \hspace{1cm} (29)

The magnitude of the deflation term summands decreases with the pole distance $|\lambda_i^{(j)} - \lambda_i^{(j)}|$. As the poles lie close to the unit circle, we can instead use the pole phase difference $\angle \left( \frac{\lambda_i^{(j)}}{\lambda_i^{(j)}} \right)$ as a proxy. The proxy distance measure is simpler to compute as we maintain a phase-sorted list of poles $\lambda^{(j)}$. The idea is then to divide the poles $\lambda^{(j)}$ into near and far pole sets, $\lambda_{\text{near } i}^{(j)}$ and $\lambda_{\text{far } i}^{(j)}$, respectively, and approximate the deflation of the less significant far poles $\lambda_{\text{far } i}^{(j)}$ by a default
term, e.g., $\lambda_{\text{far}}^{(j)}$. For symmetry, the number of near poles $N_{\text{near}}$ is assumed to be an odd number such that

$$\lambda_{\text{near}}^{(j)} = \arg \min_{\lambda \in \lambda^{(j)}, \# \lambda = N_{\text{near}}} \left| \frac{\lambda_i}{\lambda_{\text{near}}^{(j)}} - 1 \right|$$

and

$$\lambda_{\text{near}}^{(j)} = \lambda^{(j)} \setminus \lambda_{\text{near}}^{(j)}$$

where, $\#$ denotes the cardinality of a set. By sorting the system poles iterations $\lambda^{(j)}$ along pole angles, we can find the near poles $\lambda_{\text{near}}^{(j)}$ effectively. After updating the pole location $\lambda_{\text{near}}^{(j)}$ in (15), only $\lambda^{(j) \dagger}$ needs to be reordered to maintain the phase-sorted list.

It can be shown, that for equidistributed poles such as $\lambda^{(0)}$ in (26), the far deflation is (see Appendix B)

$$D_i(\lambda_{\text{far}}^{(j)}) = \frac{1}{\lambda^{(0)}} \frac{\Omega - N_{\text{near}} - 1}{2}.$$  

Thus, the total deflation may be approximated by

$$\tilde{D}_i(\lambda^{(j)}) = D_i(\lambda_{\text{near}}^{(j)}) + D_i(\lambda_{\text{far}}^{(j)})$$

if the phase of the far poles $\lambda_{\text{far}}^{(j)}$ are sufficiently uniformly distributed close to the unit circle. Occasionally, the iteration update results in a pole estimate far from the unit circle, which would compromise the far deflation approximation. In such a case, the same pole estimate is updated with a series of exact EAI steps until the pole estimate is close to the unit circle. If the initial estimates are $\lambda^{(0)}$, $\lambda^{(1)}$ may be computed only from the estimated far deflation with $N_{\text{near}} = 0$.

To establish the quality of this approximation, let us assume that there exists an upper bound $\epsilon_\text{D}$ for the approximation error of the deflation term, i.e.,

$$\epsilon_\text{D} \geq \max \left| \tilde{D}_i(\lambda^{(j)}) - D_i(\lambda^{(j)}) \right|.$$  

We can show that the error tolerance $\tau_3$ in (29) is satisfied if (see Appendix C)

$$\left| N(\lambda_{\text{near}}^{(j)})^{-1} - \tilde{D}_i(\lambda^{(j)}) \right| - \epsilon_\text{D} \geq \frac{2}{\tau_3}.$$  

In other words, if the deflation approximation is sufficiently far from the inverse Newton term, the deflation error becomes negligible. On the contrary, if the deflation term is close to the inverse Newton step, the EAI step error can be large even for small deviations in the deflation term. In case the error tolerance in (35) is not satisfied, we compute the exact deflation instead, which can be efficiently done by additionally computing the exact deflation of the far pole $D_i(\lambda_{\text{far}}^{(j)})$.

### B. Deflation Error Bound

To implement the proposed approximate deflation, we need a priori knowledge of the approximation error bound $\epsilon_\text{D}$. As $\epsilon_\text{D}$ depends on many factors such as matrix size $N$, system order $M$, feedback matrix $A$ and number of near poles $N_{\text{near}}$, in this work, it is determined experimentally from random FDNs. More precisely, for given $N$, $M$, $N_{\text{near}}$, we generate a random feedback matrix $A$ and random delays $m$ and perform the EAI method. While performing the EAI, we record each exact deflation $D_i(\lambda^{(j)})$ and approximated deflation $\tilde{D}_i(\lambda^{(j)})$.

The upper bound $\epsilon_\text{D}$ is then the maximum difference between all recorded instances of the exact and approximated deflation terms. While this heuristic approach may not yield a strict upper bound $\epsilon_\text{D}$, we have found that the EAI-AD is robust against occasional deviations.

The performance of the approximate deflation heuristics may be quantified by the number of exact deflations versus total iteration steps, i.e., the probability $P_{\text{exact}}$ that (35) is not satisfied and the exact deflation needs to be computed. Larger $N_{\text{near}}$ results in more accurate approximations and lower $P_{\text{exact}}$, while it also increases the computational load. Similarly, larger $\epsilon_\text{D}$ reduces the risk of a poor iteration step $\tilde{D}_i(\lambda^{(j)})$, but also increases $P_{\text{exact}}$. In Section V, we present practical values for these heuristic parameters.

### C. Numerical Stability

The EAI is numerically stable as the matrix inversion in (16) is only necessary if the matrix is sufficiently non-singular due to the first stopping criteria. For large delays $m$ however, the evaluation of $D_m(z)$ may become extremely large or small if $z$ is too far away from the unit circle. In particular, numerical overflow occurs for poles outside the unit circle, i.e., $|\lambda^{(j)}| > 1$ [28]. We propose a numerically stable formulation based on the reversed GCP $p_R(z) = z^m p(z^{-1})$ such that

$$p(z) = z^m p_R(z^{-1})$$

$$p'(z) = N z^{m-1} p_R(z^{-1}) - z^{-m-2} p'_R(z^{-1})$$

and (17) is therefore

$$N(z) = \frac{p(z)}{p'(z)} = \frac{1}{z^{-1} N - z^{-2} \frac{p_R(z^{-1})}{p_R(z^{-1})}}.$$  

We show in Appendix D that

$$p_R(z) = (-1)^N \det(A) \det(P_R(z)),$$  

where

$$P_R(z) = D_m(z)^{-1} A^{-1}.$$  

Similar to (17), we have

$$\frac{p_R(z)}{p_R(z)} = \text{tr}(P_R(z)^{-1} P_R'(z)).$$  

Also the first stopping criterion may be formulated stably for poles outside the unit circle. For $z \neq 0$ and $\det(A) \not= 0$, $p(z) = 0$ if and only if $p_R(z^{-1}) = 0$ if and only if $\det(P_R(z^{-1})) = 0$ if and only if $\kappa(P_R(z^{-1})) = 0$. Thus, the alternative stopping criterion is

$$\kappa(P_R(z^{-1})) < \tau_1.$$  

### V. Modal Synthesis and Evaluation

The following evaluation uses real-valued FDN parameters such that the system poles appear in complex conjugate pairs. Consequently, it is sufficient to depict the results for poles with positive phase only. In Algorithm 1, we list the pseudocode of our proposed method under evaluation.
Algorithm 1 Modal Decomposition with EAI-AD

Input: $m$, $A$, $\mathcal{H}_{\text{linear}}$, $\tau_1$, $\tau_2$, $\tau_3$

Output: $\lambda$, $\rho$

1: Compute $\epsilon_D$ from Section IV-B
2: Set $\lambda^{(0)}$ with (26)
3: Set $j \leftarrow 1$
4: repeat EAI-AD
5: for $i = 1 \ldots \mathcal{H}$ do
6: if $\kappa \left( P \left( \lambda^{(j-1)}_i \right) \right) < \tau_1$ then
7: $\lambda^{(j+1)}_i \leftarrow \lambda^{(j)}_i$
8: else
9: if $\left| \lambda^{(j)}_i \right| \leq 1$ then
10: Compute $\mathcal{N} \left( \lambda^{(j)}_i \right)$ with (17)
11:else
12: Compute $\mathcal{N} \left( \lambda^{(j)}_i \right)$ with (36)
13: Compute $\widehat{D}_i \left( \lambda^{(j)}_i \right)$ with (33)
14: if $\left| \mathcal{N} \left( \lambda^{(j)}_i \right)^{-1} - \widehat{D}_i \left( \lambda^{(j)}_i \right) \right| - \epsilon_D < \frac{2}{\tau_3}$ then
15: Compute $D_i \left( \lambda^{(j)}_i \right)$ with (18)
16: Compute $\Delta_i \left( \lambda^{(j)}_i \right)$ with (16)
17: $\lambda^{(j+1)}_i \leftarrow \lambda^{(j)}_i - \Delta_i \left( \lambda^{(j)}_i \right)$ with (15)
18: $j \leftarrow j + 1$
19: until $\forall i : \kappa \left( P \left( \lambda^{(j)}_i \right) \right) < \tau_1$ or $\kappa \left( P_R \left( 1/\lambda^{(j)}_i \right) \right) < \tau_1$
20: or $\left| \Delta_i \left( \lambda^{(j)}_i \right) \right| \leq \tau_2 \left| \lambda^{(j)}_i \right|$
21: for $i = 1 \ldots \mathcal{H}$ do Compute Residues
22: Compute $\rho_i = g(\lambda_i)$ with (20)

A. Modal Synthesis and Accuracy

A numerically accurate way to verify the modal decomposition is to synthesize each estimated mode $h_i(n)$ in time-domain as expressed in (1) and compare the sum of all modes with the impulse response $h(n)$ computed by the time-domain recursion in (3). The concept of modal synthesis and verification is depicted in Fig. 1. The error is given by the maximum difference\footnote{The maximum error is chosen as it is an upper bound for the root mean square error (RMSE) and as such a strict error measure.} between the two impulse responses, i.e.,

$$\epsilon = \max_n \left| h(n) - \sum_{i=1}^{\mathcal{H}} \hat{h}_i(n) \right|. \quad (41)$$

In the following, we evaluate FDNs with unit input and output gains, i.e., $b = 1$ and $c = 1$ such that $1 \leq \max_n h(n) \leq N$. We use double precision floating point arithmetic and the modal decomposition is regarded successful if the maximum error $\epsilon < 10^{-10}$ equivalent to -200 dB of the maximum gain of the impulse response.

B. Numerical Evaluation

For the FDN, a single EAI step in (16) can be evaluated in $O(\mathcal{H} + N^3)$: an evaluation of $P(z)$ and $P'(z)$ is merely

an evaluation of the delay matrix $D_m(z)$ in $O(N)$; a numerical matrix inversion can be performed in $O(N^3)$; and the deflation term is evaluated in $O(N)$. Thus, a full iteration from $\lambda^{(j)} \rightarrow \lambda^{(j+1)}$ can be evaluated in $O(N^2 + \mathcal{H}N^3)$. This compares favorably with the bound $O(\mathcal{H}N^3)$ of a matrix-based algorithm applied to the linearization in (11). For a high number of system poles $\mathcal{H} \gg N^3$, the complexity of computing the deflation term in (18) becomes the dominating part. The complexity of EAI-AD in (33) is similar asymptotically, however in practice, the computational complexity is reduced significantly.

Fig. 2 depicts a comparison of measured computation time with the MATLAB$^2$ functions $\text{eig}$ and $\text{roots}$ solving the direct problems (11) and (12), respectively\footnote{Matlab is a registered trademark of The MathWorks Inc. All computations were performed with Matlab R2016b on a desktop machine with an Intel Core i7 @ 3.60 GHz and 32 GB of RAM.}. The total number of delays $\mathcal{H}$ were distributed randomly among eight delay lines and the feedback matrix $A \in \mathbb{R}^{8 \times 8}$ was a random orthogonal matrix. All methods gave the correct answer with the required accuracy. For the approximate deflation, the number of near poles $\mathcal{H}_{\text{near}}$ was set to $\mathcal{H}/100$. The maximum deflation error $\epsilon_D = 10^{-3}$ was determined a priori by probing an independent set of random FDNs of similar configuration. The EAI-AD step tolerance $\tau_3$ was set to $10^{-3}$. The initialization (26) influences the performance of the EAI and EAI-AD. Nonetheless overall, we observed an average number between 4 to 5 iterations per pole.

The EAI and EAI-AD implementations utilize only standard MATLAB functions and no C-optimization which explains the relatively poor performance for small system order $\mathcal{H} < 10^3$. For high system order $\mathcal{H}$ such as $5 \cdot 10^4$, the standard EAI and EAI-AD outperform the MATLAB’s $\text{eig}$ function by a factor of more than 300 and 1300, respectively. Further, the memory requirements of the EAI and EAI-AD are only linear in $\mathcal{H}$ and

\footnote{We have omitted the evaluation of Matlab’s polynomial eigenvalue solver polyeig which performs worse than $\text{eig}$ for FDNs as it does not capture the highly sparse FDN structure.}
A. Attenuation

Attenuation filters in FDNs, as they are typically applied in artificial reverberation, aim to control the frequency-dependent reverberation time [7, 33]. As expressed in (24), all delays are extended with absorption filters $\alpha(z)$ and the feedback matrix $A$ is orthogonal or more generally unilossless [9]. We study three types of attenuation: homogeneous, near-homogeneous and inhomogeneous attenuation.

1) Homogeneous Attenuation: The attenuation filters $\alpha(z)$ are called homogeneous if there exists an attenuation-per-sample $\Gamma(z)$ such that

$$\alpha_i(z) = \Gamma(z)^m_i.$$  (42)

The attenuated delay lines can be expressed as plain delay lines with a mapped argument, i.e.,

$$D_{m}(z) \operatorname{diag}(\alpha(z)) = D_m(z \Gamma(z)^{-1}).$$  (43)

Consequently, the system poles with attenuation $\lambda_i^\Gamma$ can be related to the system poles $\lambda_i$ without attenuation by

$$\lambda_i = \lambda_i^\Gamma \Gamma(\lambda_i^\Gamma)^{-1}. $$  (44)

Now, we analyze the simple case where $\Gamma(e^{iω}) \in \mathbb{R}$. As a consequence, the mode frequencies are unaltered by the attenuation, i.e., $\angle \lambda_i = \angle \lambda_i^\Gamma$. For a unitary $A$, all unattenuated system poles $\lambda_i$ lie on the unit circle such that

$$|\lambda_i^\Gamma| = \Gamma(\lambda_i) $$  (45)

and the attenuated FDN is stable if $|\Gamma(e^{iω})| < 1$. For homogeneous attenuation, the magnitude bounds in (25) are tight. Although this simple case mainly serves a theoretical purpose, it can be realized if $\Gamma(e^{iω})$ is sufficiently smooth, i.e., attenuation filters $\alpha_i(z)$ are zero-phase FIR filters such that $z^{-m_i}\alpha_i(z)$ in (24) is causal.

2) Near-homogeneous Attenuation: Typically, the attenuation filters are implemented with relatively low order, such as one-pole filters [7], however higher order filters were proposed as well [33, 34]. The attenuation filters are designed to match the magnitude response

$$|\alpha_i(e^{iω})| \approx |\Gamma(e^{iω})^m_i|,$$  (46)

where the attenuation-per-sample is derived from a target reverberation time

$$20 \log_{10} |\Gamma(e^{iω})| = \frac{-60}{T_{60}(ω)f_s}, $$  (47)

where $f_s$ is the sampling frequency and $T_{60}(ω)$ is the time in seconds for the energy decay curve of the impulse response at frequency $ω$ to decay by 60 dB [35]. For illustration, we compute the one-pole filter according to [7] given the target reverberation time at DC $T_{60}(0)$ and Nyquist frequency $T_{60}(\pi)$. Figure 3 depicts the resulting modal decomposition for an 8-FDN with an orthogonal feedback matrix and a target reverberation time $T_{60}(0) = 2$ seconds and $T_{60}(\pi) = 0.4$ seconds. The system pole magnitudes are modified according to the target reverberation time. However, the attenuation varies especially in the transition band due to errors in the magnitude response caused by the limited filter order. The magnitude of the residues are depicted in Fig. 3b.

cubic in $N$ such that it is possible to perform modal decomposition up to $\Re = 10^6$, whereas the memory requirements for eig become prohibitive for $\Re > 5 \cdot 10^4$. For $\Re > 10^5$, more than 95% of the computation time of the EAI was spent on the deflation term, whereas it was less than 10% for the EAI-AD. For the EAI-AD, the number of exact iterations $|P|$ never exceeded 1% of the total number of iterations, proving the chosen heuristic parameters effective. The EAI-AD performs similar for small delays but outperforms the standard EAI by a factor of 100 for large $\Re$. Each EAI step is independent and only requires synchronization at every full iteration step such that the overall performance of the EAI might further be improved by parallelization.

VI. ANALYSIS OF FEEDBACK DELAY NETWORKS

We study two applications of modal decomposition in artificial reverberation: Firstly, we study the effect of attenuation filters on the poles and residues of an FDN. Secondly, we study the statistical distribution of poles and residues of random lossless FDNs.
For near-homogeneous attenuation, it can be observed that the residues with and without attenuation are rather similar. The mean difference between the residue magnitudes between corresponding poles is 0.48 dB which is small compared the overall dynamic range of 50 dB. Although the attenuation filters are not completely homogeneous, their phase component is small compared to the phase of the delays $z^{-m}$ such that the overall behavior is well approximated by (45). This suggests that studies on residues of lossless systems may translate well to results for moderately lossy systems.

3) Inhomogeneous Attenuation: While the homogeneous attenuation has perceptually desirable properties in artificial reverberation, more physically oriented FDN designs such as scattering delay networks [36] and radiance transfer [37] employ attenuation filters which are unrelated to the delay lengths but related to the boundary materials of the simulated space. Figure 4 depicts the modal decay rate of the same 8-FDN as in Fig. 3 with different attenuation filters. Instead of the delay proportional design in (46), all one-pole filters have the same target frequency response corresponding to an average delay length. As a consequence, the decay time of the neighboring modes are largely different, while the overall shape still follows the target reverberation time.

B. Statistical Distribution of Poles and Residues

We present a set of statistical analyses of lossless FDNs which rely on the proposed large-scale numerical computation of the modal decomposition and are difficult to derive by analytic methods. The statistical analysis answers a long-standing question in artificial reverberation design [38]: Why do some FDNs have an unpleasant metallic ringing despite a sufficiently high modal density? While ideal late reverberation has been characterized as Gaussian white noise [39], the metallic ringing is caused by excessive energy at few frequencies. In terms of modal decomposition, metallic ringing may be caused by either clustering of multiple poles at the ringing frequencies or largely varying energy of neighboring modes. We study the following two questions:

1) What is the distribution of the mode frequencies?

2) What is the distribution of residue magnitudes?

In the analyses, we rely on Monte Carlo simulations of randomly generated lossless FDNs.

1) Mode Frequency Distribution: The cluster number

$$C(\omega) = \# \left\{ i \mid \angle \lambda_i \in \left[ \omega - \frac{\pi}{M}, \omega + \frac{\pi}{M} \right] \right\}, \quad (48)$$

is a measure on how equally distributed the mode frequencies are. The higher the cluster number, the more poles cluster around the frequency $\omega$. In contrast, a mode gap occurs if $C(\omega) = 0$, i.e., no mode lies in this frequency interval. For perfectly equidistributed poles $C(\omega) = 1$ for all $\omega$. We evaluate the distribution of mode frequencies by computing the probability of the integer cluster number $\eta$, i.e., $P(C(\omega) = \eta)$. The random 8-FDNs have delays between 50 and 1000 samples and an orthogonal feedback matrix. The probabilities are averaged over 100 random instances each.

In Table I, the probability of cluster numbers for randomly generated FDNs are compared to cluster numbers of a pseudo-uniform random number generator with equal sampling size. The discrepancy of the cluster number from an equidistribution is relatively low for the FDN modes compared to the random number generator. In fact for FDNs, it is very rare to find an interval of width $2\pi/\Omega$ with more than two modes. In stark contrast, acoustic mode density of physical spaces increase quadratically with frequency [12].

2) Residue Magnitude Distribution: In Section II-D, we have presented the computation of the mode residues for a
given set of system poles. Figure 5 depicts the magnitude histogram of the total and undriven residues as well as the input-output drives for a random 8-FDN. The input-output drives are comprised of all individual input-output combinations, i.e., \( \text{adj}(P(\lambda_i)) \) in (22). The total residues \( \rho(\lambda_i) \) result from unit input and output gains, i.e., \( b = 1 \) and \( c = 1 \), or in other words, \( \rho(\lambda_i) = \rho^2(\lambda_i)(1 - \text{adj}(P(\lambda_i))1) \). The magnitude distributions of the inverse undriven residues \( 1/\rho^2(\lambda_i) \), the total residues \( \rho(\lambda_i) \) and input-output drives \( \text{adj}(P(\lambda_i)) \) all resemble log-Rayleigh distributions [40]. However, just by altering the feedback matrix \( A \), it is possible to encounter various other distributions of the residue magnitude. Figure 6 depicts the residue magnitude distribution of four selected orthogonal feedback matrices.

3) Discussion: For randomly generated lossless FDNs, the mode frequencies are nearly equidistributed such that every frequency band has energy contributions from a similar number of modes. On the other hand, the wide distribution of the residue magnitudes especially around the upper limit suggests that a small number of poles contribute a large portion of the impulse response energy. In Fig. 7, we have synthesized the impulse responses corresponding to the modal decomposition of the FDNs in Fig. 6. We first sort the modes along the residue magnitude in descending order, i.e., \( |\rho_i| \geq |\rho_1| \) for \( i < l \). The signal power error (SPE) of the time-domain impulse responses of length \( L \) are

\[
\epsilon_{\text{SPE}} = \frac{\left\| h(n) - \sum_{i=1}^{\rho} \tilde{h}_i(n) \right\|^2}{L},
\]

(49)

where the impulse responses were 5 seconds at 48 kHz sample rate long, i.e., \( L = 240000 \), and the power of \( h(n) \) was normalized, i.e., \( \|h(n)\|^2/L = 1 \). For feedback matrix \( A_1 \), 20% of the highest energy modes result in \( \epsilon_{\text{SPE}} = 0.47 \), whereas the \( \epsilon_{\text{SPE}} = 0.03 \) for \( A_4 \). In the context of artificial reverberation, the high-energy modes dominate the frequency spectrum such that the audible modal density is considerably lower than theoretic modal density, i.e., the number of modes per frequency. For illustration, we have synthesized audio examples from the four instances depicted in Fig. 6 and provided them online\(^4\). The impulse responses from these lossless FDNs do not decay in energy over time and therefore do not represent examples for artificial reverberation. Nonetheless, as demonstrated earlier such a lossless FDN can be converted into a lossy FDN by introducing additional attenuation filters. The impulse response from feedback matrix \( A_1 \) resembles the sound of white noise, while the impulse response for feedback matrix \( A_4 \) has strong coloring resonances as we expect from residue magnitude distribution in Fig. 6 and Fig. 7.

The residue distribution may be optimized in two steps: Firstly, optimization of the undriven residues by choosing delays \( m \) and feedback matrix \( A \). Secondly, optimization of the total residue by choosing the input and output gains, \( b \) and \( c \). While the first step is a non-linear process which requires further research, the second step may be readily solved by linear least square fitting.

\(^4\)www.sebastianjiroschlecht.com/publication/schlecht-2019-uj/

VII. Conclusion

We presented a numerically efficient technique for modal decomposition of the FDN. Standard methods such as eigenvalue decomposition of the linearized system and polynomial root finding methods applied to the characteristic polynomial require significant computational resources when the system order is large. The proposed method applies the Ehrlich-Aberth Iteration (EAI) to the polynomial matrix formulation of the FDN. Further we proposed novel bounds on the pole locations and an efficient approximate deflation technique based on the estimation of far poles (EAI-AD). For high system order such as \( 5 \times 10^4 \), the EAI and EAI-AD outperform the MATLAB’s \text{eig} function by a factor of more than 300 and 1300, respectively. The EAI-AD was able to give reliable results up to a system order of 1 million. The modal decomposition was applied to FDNs in the context of artificial reverberation. Three types of attenuation were studied: homogeneous, near-homogeneous and inhomogeneous. The potential for explicit analysis of the pole and residues was demonstrated for attenuation filter design. Statistical analysis showed that for randomly generated FDNs, the mode frequencies are nearly equidistributed and the residue magnitudes follow a log-Rayleigh distribution. This analysis suggests that relatively few modes are contributing a large portion of the late reverberation energy. Overall, the connection to FDN parameters, in particular, the feedback matrix, facilitates future research to improve the parametrization and the quality of artificial reverberation.

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Appendix A

Bounds of Pole Magnitudes

We present lower and upper bounds on the pole magnitudes \( |A| \) of an FDN. The bounds are based on the generalization of Rouché’s theorem to matrix polynomials.
Let $S(z)$ and $Q(z)$ be matrix polynomials and let $\gamma$ be a simple closed Jordan curve. If $S(z) A - Q(z) X\geq 0 \text{ for } |z| = r \bigcup X - Y \text{ is positive definite, then the polynomials } \det(S(z)) \text{ and } \det(S(z) + Q(z)) \text{ have the same number of roots in the open set bounded by } \gamma.$

A simple choice of the Jordan curve $\gamma$ is the circle with radius $r$, i.e., $|z| = r$. An immediate consequence of the above theorem applied to the polynomial $P(z)$ of (7) with $S(z) = -A$ and $Q(z) = D(m) - 1$ is (21): If

$$A^H A \succ D_m(r^{-2}) \tag{51}$$

which in turn is equivalent to [43]

$$\rho \left( (A^H A)^{-1} D_m(r^{-2}) \right) \leq 1, \tag{52}$$

where $\rho(X)$ denotes the spectral radius of a matrix $X$. Using properties of the spectral norm [43] we can give an upper bound on this expression by

$$\rho \left( (A^H A)^{-1} D_m(r^{-2}) \right) \leq \left\| (A^H A)^{-1} D_m(r^{-2}) \right\|_2 \leq \left\| A^{-1} \right\|_2 \left\| A^{-H} \right\|_2 \left\| D_m(r^{-2}) \right\|_2 \tag{53}$$

Thus, the criterion (52) is satisfied if

$$r^2 \min m = \left\| D_m(r^{-2}) \right\|_2 \leq \left\| A^{-1} \right\|_2^2 = \min \sigma(A)^2. \tag{54}$$

Therefore with (50), the pole magnitude lower bound may be given as

$$\min |\lambda| \geq \min \sigma(A)^{1/min m}. \tag{55}$$

Analogously, applying the same arguments to the reversed matrix polynomial $P_R(z^{-1})$ yields an upper bound

$$\max \{|\lambda| \leq \max \sigma(A)^{1/\max m}. \tag{56}$$

For additional attenuation filters $\alpha(z)$ as in (24) and a unitary feedback matrix $A$, i.e., $A^H A = I$, these bounds can be tightened further. Rouche’s criterion with $S(z) = -A$ and $Q(z) = D_m(z)^{-1} \alpha(z)^{-1}$ in Theorem 1 gives the condition

$$|\alpha(z)| \geq |z|^m. \tag{57}$$

Let $\gamma$ be a simple closed curve with

$$\gamma(\omega) = \min \left( |\alpha(e^{i\omega})|^{1/m} \right) e^{i\omega}. \tag{58}$$

For large delays $m$, we have $\alpha(\gamma(\omega)) \approx \alpha(e^{i\omega})$ such that (57) is satisfied for $z = \gamma(\angle z)$. Thus, the lower bound of the pole magnitude is

$$|\lambda_i| \geq \min \left( |\alpha(e^{i\angle \lambda_i})|^{1/m} \right). \tag{59}$$

The corresponding upper bound may be derived similar to (56) by applying Rouche’s theorem to $P_R(z^{-1})$:

$$|\lambda_i| \leq \max \left( |\alpha(e^{i\angle \lambda_i})|^{1/m} \right). \tag{60}$$

These bounds are tight for a diagonal matrix $A$.

**APPENDIX B**

**Far Deflation Estimation**

We are given the equidistributed poles $\lambda^{(0)}$ as defined in (26) and an even number of near poles $\eta_{\text{near}}$. We compute the far deflation for pole $\lambda^{(0)}$. First we state a useful identity. For any real $x$,

$$\frac{1}{1-e^{ix}} + \frac{1}{1-e^{-ix}} = 1. \tag{61}$$

The total deflation is

$$D_t(\lambda^{(0)}) = \sum_{l=1}^{\infty} \frac{1}{\lambda^{(0)}_l - \lambda^{(0)}_j} \tag{62}$$

Similarly, as each conjugate pair of poles contribute equally to the deflation, the far deflation is

$$D_t(\lambda^{(0)}_{\text{far, i}}) = \frac{1}{\lambda^{(0)}_i} \frac{\eta - \eta_{\text{near}} - 1}{2}. \tag{63}$$

**APPENDIX C**

**Deflation Error**

We show that inequality (35) is sufficient for inequality (29). For the sake of brevity, we omit the pole arguments in the following. Given the deflation approximation $D_t$, which satisfy (35), we obtain

$$\frac{\tau_3}{2} \geq \frac{1}{N-1-D_t - \epsilon_D} \geq 0. \tag{64}$$
We show that (63) satisfies EAI step error tolerance (29). Because $\epsilon_D \geq 0$, it is
\[ \frac{\tau_3}{2} \geq \frac{1}{|N^{-1} - D_i|} \geq 0. \] (64)
Further, as $|D_i - \tilde{D}_i| - \epsilon_D \leq 0$,
\[ 1 \left| \frac{N^{-1} - D_i}{|N^{-1} - D_i| + |D_i - \tilde{D}_i| - \epsilon_D} \right| \leq \frac{1}{N^{-1} - D_i} \leq \frac{\tau_3}{2} \]
Eventually, we can show that
\[ |\Delta_i^{(j)} - \tilde{\Delta}_i^{(j)}| \leq \left| \frac{N^{-1} - D_i}{|N^{-1} - D_i| + |D_i - \tilde{D}_i| - \epsilon_D} \right| \leq \frac{\tau_3}{2}. \]

**APPENDIX D**

**REVERSED CHARACTERISTIC POLYNOMIAL**

In [18], the authors showed that the coefficient $c_i$ of the GCP in (12) are
\[ c_i = \begin{cases} \sum_{T \in I_i}(-1)^{|I|} \det(A(I^c)), & \text{for } I_i \neq \emptyset \\ 0, & \text{otherwise} \end{cases}, \] (65)
where $\langle N \rangle = \{1, 2, \ldots, N\}$ and $I_i = \{I \subseteq \langle N \rangle \mid \sum_{i \in I} m_i = i\}$ is the set of index sets $I$ such that the delays with indices in $I$ sum up to $i$. Further, $I^c$ is the relative complement in $\langle N \rangle$, i.e., $I^c = \langle N \rangle \setminus I$ and $A(I^c)$ indicates the submatrix with rows and column indices in $I^c$. With Jacoby’s identity for invertible matrices $A$
\[ \det A^{-1}(I) = \frac{\det A(I^c)}{\det A}, \] (66)
and
\[ I \in I_{\Omega_i} \text{ if and only if } I^c \in I_i \] (67)
we have for $I_i \neq \emptyset$
\[ c_{\Omega_i} = (-1)^N \det(A) \sum_{I \in I_i} (-1)^{|I|} \det(A(I^c)). \]
Consequently,
\[ p_R(z) = (-1)^N \det(A) \det(P_R(z)). \] (68)

**REFERENCES**


